

AOCC User Guide

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Revision History

| Date | Revision | Description |
|---------------|----------|---|
| November 2022 | 4.0 | Added Chapter 7.Included release specific updates and reorganized content. |
| December 2021 | 3.2 | Updated Chapter 5 with Flang sanitization command. Updated the Chapter 4 with (NO)FREEFORM pragma related information. |
| July 2021 | 3.1 | Incremental updates and some major/general edits. |
| March 2021 | 3.0 | Initial version. |

Chapter 1 Introduction

The AMD Optimizing C/C++ and Fortran Compiler (AOCC) is highly optimized for x86 targets, especially for AMD "Zen"-based processors. This guide describes how to use AOCC.

AOCC 4.0 is based on the LLVMTM 14.0.6 compiler infrastructure (llvm.org, 24 June 2022) and includes bug fixes and support for other new features. For more information, refer to AOCC 4.0 release notes.

Chapter 2 Programming Language Support

AOCC is a high-performance x86 CPU compiler for C, C++, and Fortran programming languages. It supports target-dependent (x86 targets especially AMD processors) and target-independent optimizations.

AOCC leverages LLVM Clang for the compiler and driver for C and C++ programs. Flang is the compiler and driver for Fortran programs. Clang supports 32 and 64-bit targets while Flang supports only 64-bit targets.

2.1 C, C++, AND FORTRAN Programming Languages

AOCC Clang and AOCC Flang support the preprocessing, parsing, optimization, code generation, assembly, and linking. Using these drivers, overall execution of other tools, such as the compiler, assembler, and linker can be controlled depending on which high-level mode setting is passed. While Clang and Flang are highly integrated, it is important to understand the stages of compilation. These stages are executed in the following sequence:

1. **Driver**

Clang is not just a C and C++ front-end that compiles the program to LLVM intermediate representation (IR). Clang is also the driver that ensures use of the required LLVM optimization passes and targets code generation to generating the binaries.

Similar to Clang, Flang is the Fortran front-end compiler and consists of the following two components:

- flang1: Invoked by the front-end driver responsible for transforming the Fortran programs into tokens. The parser transforms these tokens into Abstract Syntax Tree (AST). The AST is then transformed into canonical form that is used to generate the ILM code.
- flang2: Uses the ILM code from flang1 and transforms it into ILI that is then optimized by the internal optimizer. The optimized ILI is then transformed into LLVM IR. Then, the front-end driver transfers this LLVM IR to LLVM optimizer for optimization and target code generation.

For simplicity, you can use these Clang and Flang as an end-to-end driver. However, for advanced compilation, you can manually execute each compilation phase.

2. Preprocessing

This stage handles the tokenization of the input source file, macro expansion, #include expansion, and handling of the other preprocessor directives. The output of this stage is typically called a .*i* (for C), .*ii* (for C++), or .*i* (for Fortran) file.

3. Parsing and Semantic Analysis

This stage parses the input file, translating the preprocessor tokens into a parse tree. When in the form of a parse tree, it applies semantic analysis to the compute types for expressions and to determine whether the code is well-formed. This stage is responsible for generating most of the

compiler warnings and parse errors. The output of this stage from Clang is an Abstract Syntax Tree (AST). Using Flang, Flang1 will be invoked to transform the program tokens into AST, then into canonical form- that is used to generate ILM code.

4. LLVM IR Code Generation

In Clang, this stage translates an AST into a Low-level Intermediate Code (LLVM IR)

Using Flang, Flang2 takes up the ILM code generated by Flang1 and transforms it into ILI, which is then optimized by the internal optimizer and then transformed into LLVM IR.

5. AOCC Optimizer

This phase is responsible for optimizing the generated LLVM IR and handling the target-specific code generation. The output of this stage is typically called a .s or an assembly file.

6. Machine Code Generation

This phase performs the target specific code generation from the optimized LLVM IR. The output of this stage is typically called a .s or assembly file. Clang and Flang also support the use of an integrated assembler from which the code generator produces object files directly. This avoids the overhead of generating the .s file and then calling the target assembler.

7. Assembler

This stage runs the target assembler to translate the output of the compiler into a target object file. The output of this stage is typically called a .o or object file.

8. Linker

This stage runs the target linker to merge multiple object files into an executable or dynamic library. The output of this stage is typically called an *a.out*, .*dylib*, or .*so* file.

2.2 Compatibility/Conformance to Standards

AOCC supports the following language and debugging standards:

- C: C99, C11, and C17 standards
- C++: C++ 98, C++ 03, C++ 11, C++ 17, and C++ 20 standard¹
- Fortran: F77, F90, F95, F2003, and F2008 standards²
- OMP 4.5 and OMP 5.0 standards for C/C++ programming³
- OMP 4.5 standards for Fortran programming
- DWARF 5 standards for C, C++, and Fortran debuggability

^{1.} Refer https://clang.llvm.org/cxx status.html#cxx20 for C++ 20 conformance.

^{2.} AOCC does not support F2008 coarrays.

^{3.} Refer https://releases.llvm.org/13.0.0/tools/clang/docs/OpenMPSupport.html#openmp-implementation-details for OpenMP 5.0 conformance.

2.2.1 Support for Annex F (IEEE-754/IEC 559) of C99/C11

Clang compiler does not support IEC 559 math functionality. Clang does not control and honor the definition of __STDC_IEC_559__ macro. Under specific options, such as -Ofast and -ffast-math, Clang will enable a range of optimizations that provide faster mathematical operations that may not conform to the IEEE-754 specifications. The macro __STDC_IEC_559__ value may be defined but ignored when these faster optimizations are enabled.

2.2.2 IEEE-754 Support

The Flang compiler does not conform to IEEE-754 specifications when -ofast or -ffast-math options are specified. The compiler will enable a range of optimizations that provide faster mathematical operations under -ofast and -ffast-math mode of compilation.

Note: AOCC Flang extends the GitHub version (https://github.com/flang-compiler/flang.git) with enhancements and stability.

Chapter 3 Working with AOCC

3.1 Installing On Linux

3.1.1 Prerequisites

The following software packages must be installed prior to the AOCC installation:

Table 1. Prerequisites

| Package Name | Version(s) | Notes |
|----------------|----------------|--|
| libstdc++ | 6 or later | GNU Standard C++ Library V3 |
| libncurses-dev | 5.9 or later | Provides libtinfo, which is a low level terminfo library |
| zlib | 1.2.7 or later | Compression library |
| Libxml2 | 2 or later | Parses the XML documents |
| libquadmath | 4.8 or later | GCC Quad-Precision Math Library |
| python | 3.x | Python library |

Notes:

- 1. For a better performance, it is recommended to use the latest versions of Glibc and Binutils.
- 2. AOCC compiler binaries are suitable only for the Linux[®] systems having Glibc version 2.17 or later.

3.1.2 Installation

Note: This installation does not require root or sudo permission.

To install *aocc-compiler-<ver>.tar*, execute the following commands:

- 1. cd <compdir>
- 2. tar -xvf aocc-compiler-<ver>.tar
- 3. cd aocc-compiler-<ver>
- 4. bash install.sh

It will install the compiler and display the AOCC setup instructions.

5. source <compdir>/setenv_AOCC.sh

This will setup the shell environment for using AOCC C, C++, and Fortran compiler where the command is executed.

You must ensure the following:

- Run the bash command <compdir>/aocc-compiler-<ver>/AOCC-prerequisites-check.sh to check if you have all the prerequisites and your shell environment is configured correctly.
 - If there are failing checks, correct them (repeat any of the above steps that you may have missed)
 and run prerequisites check.sh again.
 - Repeat until AOCC-prerequisites-check.sh displays Check:PASSED.

Note: You could proceed if the packages mentioned in the warnings during the failing checks are not required for your run.

• The compiler is installed and your environment is configured to the current release of AOCC. At any point, you can execute the command source <compdir>/setenv_AOCC.sh to set the environment variables for the installed compiler.

3.1.3 SPACK Support

Notes:

- 1. SPACK support is available starting from AOCC 2.2.
- 2. In the following steps, <Version Number> in aocc@<Version Number> implies the AOCC version. For example, if you are installing AOCC 4.0.0, you must use aocc@4.0.0.

Installing AOCC in SPACK

Complete the following steps to installing AOCC compiler in SPACK:

1. Install AOCC:

```
$ spack install -v aocc@<Version Number> +license-agreed
```

2. Add AOCC to the SPACK compiler list:

```
$ spack cd -i aocc@<Version Number>
$ spack compiler add $PWD
$ spack cd -i aocc@<Version Number>
$ spack compiler add $PWD
```

3. List all the available compilers:

\$ spack compilers

Uninstalling AOCC Compiler from SPACK

1. Uninstall AOCC:

\$ spack uninstall aocc@<Version Number>

2. Remove the compiler from the *compiler.yaml* file:

```
$ spack compiler remove aocc@<Version Number>
```

For more details on AOCC in SPACK, refer AMD Developer Central (https://developer.amd.com/spack/amd-optimized-c-cpp-compiler/).

3.1.4 Upgrading AMD LibM (ALM)

This is required only when you are upgrading AMD LibM from the AMD portal (https://developer.amd.com/amd-aocl/).

Complete the following steps to perform an upgrade:

- 1. Extract the latest AMD LibM package.
- 2. Overwrite *aocc-compiler-<ver>/lib/libalm.so* and *aocc-compiler-<ver>/lib/libalm.a* with the latest versions of *libalm.so* and *libalm.a* respectively.
- 3. Similarly, overwrite *aocc-compiler-<ver>/include/amdlibm.h* and *amdlibm_vec.h* with the latest versions of *amdlibm.h* and *amdlibm vec.h* respectively.

3.1.5 Supported Operating Systems (OS)

The following OS are supported in this release:

- RHEL 8.6 and 9.0
- SLES 15 SP3
- Ubuntu 22.04 LTS
- CentOS 8
- Other Linux flavors/versions with glibc 2.17 or higher

3.1.6 Known Issues and Limitations

This release has the following known issues and limitations:

- AOCC binaries can run optimally only on Linux systems having glibc version 2.17 or later.
- Currently, Flang supports only 64-bit targets.

3.2 Invoking AOCC

To set the required environment before invoking compiler driver:

\$ source <compdir>/setenv_AOCC.sh

3.2.1 AOCC Optimizer

AOCC includes many optimizations for independent and dependent targets. Specific optimizations are made default when you use an optimization level *O3* and above. You can read more about these in the command line option section. Some optimizations need a whole program analysis and are enabled under Link Time Optimization (LTO) using -flto. The AOCC preferred linker is LLD. Refer the section LLD Linker for using LLD in the complier driver.

3.2.2 Using the Compiler

3.2.2.1 Clang and Clang++

To build and run a C or C++ program, execute the following commands:

```
$ clang [command line flags] xyz.c -o xyz.out
$ ./xyz.out
$ clang++ [command line flags] xyz.cpp -o xyz.out
$ ./xyz.out
```

3.2.2.2 Flang

To build and run Fortran programs, execute the following commands:

```
$ flang [command line flags] xyz.f90 -o xyz.out
$ ./xyz.out
```

3.2.2.3 LLD Linker

To use an LLD linker, execute the following commands:

```
$ clang [command line flags] -fuse-ld=lld xyz.c abc.c -o xyz.out [here -fuse-ld=lld is optional
as this option is default]
$ ./xyz.out
```

3.2.3 Libraries

Some applications may perform better using the AMD Optimizing CPU Libraries (AOCL). AOCC will work seamlessly with these libraries. It is recommended that you evaluate these libraries while building your application with AOCC. For more information on AMD Optimizing CPU Libraries (AOCL), refer AMD Developer Central (https://developer.amd.com/amd-aocl/).

3.2.3.1 Configuring Library Path

Execute the following commands to configure the library path:

- For 64-bit Library:
 - export LD_LIBRARY_PATH=<compdir>/aocc-compiler-<ver>/lib:\$LD_LIBRARY_PATH
- For 32-bit Library:

```
export LD_LIBRARY_PATH=<compdir>/aocc-compiler-<ver>/lib32:$LD_LIBRARY_PATH
```

For other AMD optimizing CPU libraries
 export LD_LIBRARY_PATH=<Path to AMD optimizing CPU Libraries>:\$LD_LIBRARY_PATH

3.2.3.2 Linking AMD Library

AMDLIBM

Execute the following command to link AMDLIBM with the linker:

\$ clang [command line flags] xyz.c -L<compdir>/aocc-compiler-<ver>/lib -lamdlibm -lm -o xyz.out

Execute the following command to link other AMD optimizing CPU libraries with linker:

\$ clang [command line flags] xyz.c -L<Path to AMD optimizing CPU Libraries> -l<library name> -o
xyz.out

AMD Vector Library

Execute one of the following commands to link the vector library calls from AOCC:

```
$ clang [command line flags] xyz.c -lamdlibm -fveclib=AMDLIBM -lm -o xyz.out
$ clang [command line flags] xyz.c -lamdlibm -mllvm -vector-library=AMDLIBM -lm -o xyz.out
```

AMD FASTMATHLibrary

-fsclrlib=AMDLIBM with -Ofast enables the use of the fastmath functions from amdlibm. Linking may fail with undefined symbols if the correct library is not used for linking.

Execute the following command to link AMDLIBMFAST with the linker:

```
$clang -Ofast [other command line flags] xyz.c -lamdlibm -fsclrlib=AMDLIBM -lamdlibmfast -lm -
o xyz.out
```

amdalloc

amdalloc is AMD's memory allocator based on jemalloc. Two versions of amdalloc, namely amdalloc and amdalloc-ext are available in the AOCC binary package. amdalloc-ext is tuned for workloads that use smaller allocation sizes.

Execute the following command to link amdalloc with the linker:

```
$ clang [command line flags] xyz.c -L<Path to AMD optimizing CPU Libraries> -lamdalloc -o
xyz.out
```

Chapter 4 Using Pragma Directives

4.1 Flang

Following are the pragma directives specific only to Flang:

4.1.1 NOINLINE

This directive instructs the compiler not to inline the specified routine.

!DIR\$ NOINLINE

To use this directive, compiler optimization level should be in -O0 to -O3. The NOINLINE directive overrides the compiler options -finline-functions and -fno-inline-functions.

Example:

```
!DIR$ NOINLINE

SUBROUTINE func_noinline

INTEGER :: i

do i = 0, 5

WRITE(*, *) "Hello World"

end do

END SUBROUTINE func_noinline

PROGRAM test_inline

IMPLICIT NONE

call func_noinline

END PROGRAM test inline

END PROGRAM test inline
```

4.1.2 FORCEINLINE

This directive instructs compiler to always inline the specified routine

!DIR\$ FORCEINLINE

To use this directive, compiler optimization level should be in -O0 to -O3. The FORCEINLINE directive overrides the compiler options -finline-functions and -fno-inline-functions.

Example:

```
!DIR$ FORCEINLINE

SUBROUTINE func_forceinline

INTEGER :: i

do i = 0, 5

WRITE(*, *) "Hello World"

end do

END SUBROUTINE func_forceinline

PROGRAM test_inline

IMPLICIT NONE

call func_forceinline

END PROGRAM test_inline

END PROGRAM test_inline
```

4.1.3 UNROLL

This directive instructs the compiler about the number of times the loop should be unrolled. !DIR\$ UNROLL [(n)]

- **n** optional parameter, integer constant ranges from 1 512
- When **n** equals 0, compiler will decide if unrolling should happen or not

To use this directive, compiler optimization level should be -O1 or above.

If **n** is specified, the optimizer unrolls the loop by n times.

If **n** is not specified or out of range, the optimizer unrolls the loop based on profitability.

Example:

```
Example 1:
subroutine func1(a, b)
 integer :: m = 10
 integer :: i, a(m), b(m)
  !dir$ unroll
  do i = 1, m
   b(i) = a(i) + 1
  end do
end subroutine func1
Example 2:
subroutine func2(m, a, b)
 integer :: i, m, a(m), b(m)
  !dir$ unroll(4)
 do i = 1, m
    b(i) = a(i) + 1
  end do
end subroutine func2
```

4.1.4 **NOUNROLL**

This directive disables unroll of the loop before which it has been used and is the opposite of UNROLL.

!DIR\$ NOUNROLL

Example:

```
subroutine func1(a, b)
  integer :: m = 10
  integer :: i, a(m), b(m)

!dir$ nounroll
  do i = 1, m
     b(i) = a(i) + 1
  end do
end subroutine func1
```

4.1.5 PREFETCH

This directive is used to insert a hint in the code generator to prefetch instruction for memory references, wherever supported. This allows a better performance in the characteristics of the code. For more information, refer LLVM documentation (https://llvm.org/docs/LangRef.html#llvm-prefetch-intrinsic).

!\$MEM PREFETCH

Constraints: To enable this directive, compiler optimization level should be in -O0 to -O3.

Example:

```
subroutine prefetch_dir(a1, a2)
  integer :: a1(4096)
  integer :: a2(4096)

do i = 128, (4096 - 128)
  !$mem prefetch a1, a2(i + 256)
   a1(i) = a2(i - 127) + a2(i + 127)
  end do
end subroutine prefetch_dir
```

4.1.6 Vectorization Pragmas

The compiler directives to control the loop vectorizations are:

- !DIR\$ VECTOR
- !DIR\$ NOVECTOR
- !DIR\$ VECTOR ALWAYS

For these pragmas to be applicable, the optimization levels should be in -O1 to -O3.

-Menable-vectorize-pragmas introduced in AOCC 3.1 (to control the vectorization pragmas) is deprecated in AOCC 3.2.

4.1.7 FREEFORM/NOFREEFORM

FREEFORM makes the compiler compile the source in free-form format, while NOFREEFORM directs to compile in a fixed-form format. They are applied to the rest of the file in which they are mentioned, but the directive is reverted if the compiler finds the opposite directive in the same file.

Example:

Chapter 5 Command-line Options

5.1 Clang and Flang Options

5.1.1 Target Selection

Following is the list of all the target selection options:

• -march=<cpu>

Enables Clang/Flang to generate code for a specific processor of a family. For example, if you specify -march=znver1, the compiler can generate instructions that are valid on AMD 17h and later processors.

-march=znver1

Use this architecture option for enabling the best code generation and tuning for AMD "Zen"-based x86 architecture. All the x86 AMD "Zen" ISA and associated intrinsic are supported.

-march=znver2

Use this architecture option for enabling the best code generation and tuning for AMD "Zen2"-based x86 architecture. All x86 AMD "Zen2" ISA and associated intrinsic are supported.

• -march=znver3

Use this architecture option for enabling best code generation and tuning for AMD "Zen3"-based x86 architecture. All x86 AMD "Zen3" ISA and associated intrinsic are supported.

-march=znver4

Use this architecture option for enabling best code generation and tuning for AMD "Zen4"-based x86 architecture. All x86 AMD "Zen4" ISA and associated intrinsic are supported.

5.1.2 Driver

-mllvm <options>

Need to provide -mllvm, so that, the option can pass through the compiler front end and is applied on the optimizer where this optimization is implemented.

For example, -mllvm -enable-strided-vectorization

-fuse-ld=lld

To invoke Ild linker from compiler driver as it is the preferred linker.

Note: For more information on the Clang options, refer Clang Documentation (https://releases.llvm.org/14.0.0/tools/clang/docs/ClangCommandLineReference.html).

5.2 Flang Options

For a list of compiler options, use the following commands:

```
$flang -help
$flang --help-hidden
```

The Flang compiler supports all the Clang compiler options (http://clang.llvm.org/docs/CommandGuide/clang.html) and the following Flang-specific compiler options:

• -Kieee

It is enabled by default from AOCC 2.2.0.

It instructs the compiler to conform to the IEEE-754 specifications. The compiler will perform floating-point operations in strict conformance with the IEEE 754 standard. Some optimizations are disabled when this option is specified.

• -no-flang-libs

Do not link against Flang libraries.

-mp

Enable OpenMP and link with OpenMP library libomp.

nomp

Do not link with OpenMP library *libomp*.

• -Mbackslash

Treat backslash character like a c-style escape character.

-Mnobackslash

Treat backslash like any other character.

-Mbyteswapio

Swap byte-order for unformatted input/output.

-Mfixed

Assume fixed-format source.

-Mextend

Allow source lines up to 132 characters.

-Mfreeform

Assume free-format source.

-Mpreprocess

Run preprocessor for Fortran files.

-Mstandard

Check standard conformance.

-Msave

Assume all variables have SAVE attribute.

• -module

Path to module file (-I also works).

• -Mallocatable=95

Select Fortran 95 semantics for assignments to allocatable objects (default).

-Mallocatable=03

Select Fortran 03 semantics for assignments to allocatable objects.

-static-flang-libs

Link using static Flang libraries.

• -M[no]daz

Treat denormalized numbers as zero.

• -M[no]flushz

Set SSE to flush-to-zero mode.

-Mcache_align

Align large objects on cache-line boundaries.

-M[no]fprelaxed

This option is ignored.

-fdefault-integer-8

Treat INTEGER and LOGICAL as INTEGER*8 and LOGICAL*8.

-fdefault-real-8

Treat REAL as REAL*8.

• -i8

Treat INTEGER and LOGICAL as INTEGER*8 and LOGICAL*8.

• -r8

Treat REAL as REAL*8.

• -fno-fortran-main

Do not link in Fortran main.

• -Mrecursive

Allocate local variables on the stack; thus, allowing recursion. SAVEd, data-initialized, or namelist members are always allocated statically, regardless of the setting of this switch.

5.3 Code Generation and Optimization Options

Both Clang and Flang rely on AOCC optimizer and code generator stages to transform the LLVM IR and generate the best code for the target x86 platform.

Following is the list of optimization options categorized by type:

Optimization Level Options

-O0 (default)

No optimization: this level compiles the fastest and generates the most debuggable code.

-O1

Between the levels -O0 and -O2.

-O2

Enables most optimizations.

-O3

Enables all optimizations, which take longer to perform or may generate larger code (in an attempt to make the program run faster).

The -O3 level in AOCC has more optimizations than the base LLVM version. These optimizations include improved handling of indirect calls and advanced vectorization.

-Ofast

Enables all the optimizations from -O3 along with other aggressive optimizations that may violate strict compliance with language standards.

The -Ofast level in AOCC has more optimizations than the base LLVM version. These optimizations include partial unswitching, improvements to inlining, and unrolling.

• -Os

Similar to the level -O2, but with extra optimizations to reduce the code size.

-Oz

Similar to the level -Os (and thus, -O2), but reduces the code size further.

• -O

Equivalent to the level -O1.

O4 and higher

Equivalent to the level -O3.

-zopt

Enables the subset of scalar, vector, and loop transformations including improved variants of loop invariant code motion, SLP and loop vectorizations, loop-fusion, loop-interchange, loop-unswitch, loop tiling, and loop distribution

Vector Optimization Options

• -enable-strided-vectorization

Enables strided memory vectorization as an enhancement to the interleaved vectorization framework present in LLVM. It enables the effective use of gather/scatter instruction patterns. This option must be used along with the interleave vectorization option.

Usage: -mllvm -enable-strided-vectorization

-global-vectorize-slp={true,false}

Vectorizes the straight-line code inside a basic block with data reordering vector operations. This option is set to **false** by default.

Usage: -mllvm -global-vectorize-slp={true, false}

region-vectorize

Enables vectorization on certain loops with complex control flow which the normal vectorizer cannot handle.

Usage: -Wl,-mllvm -Wl,-region-vectorize

-enable-loop-vectorization-with-conditions

Enables efficient vectorization of loops with conditions by conditionally executing the vector instructions as opposed to flattening the loop body and vectorizing. The vectorized code uses vector versions of compare instructions to guard the instructions in the loop body and uses masked instructions to guard against unsafe memory operations. By default, this option is enabled under O3.

 ${f Usage:}$ -mllvm -enable-loop-vectorization-with-conditions

-legalize-vector-library-calls

Splits up the unsupported higher vector factor version of vector library calls into supported vector factor version of vector library calls for enabling vectorization. By default, this option is enabled under O3.

Usage: -mllvm -legalize-vector-library-calls

-vectorize-non-contiguous-memory-aggressively

Enables vectorization involving non-contiguous memory locations by generating multiple loads/stores and inserts. By default, this option is enabled under O3.

Usage: -mllvm -vectorize-non-contiguous-memory-aggressively

-fvector-transform

Enables the subset of vector transformations including improved variants of SLP and loop vectorization. It is enabled under -zopt.

Usage: -fvector-transform

• -fepilog-vectorization-of-inductions

Enables epilog vectorization of loops that require vectorization of loop induction variables.

Usage: -fepilog-vectorization-of-inductions

Loop Optimization Options

• -enable-partial-unswitch

Enables partial loop un-switching, which is an enhancement to the existing loop unswitching optimization in LLVM. Partial loop un-switching hoists a condition inside a loop from a path for which the execution condition remains invariant, whereas the original loop un-switching works for a condition that is completely loop invariant. The condition inside the loop gets hoisted out from the invariant path and original loop is retained for the path where condition is variant.

Usage: -mllvm -enable-partial-unswitch

-aggressive-loop-unswitch

Enables aggressive loop unswitching heuristic (including -mllvm -enable-partial-unswitch) based on the usage of the branch conditional values. Loop unswitching leads to code-bloat. Code-bloat can be minimized if the hoisted condition is executed more often. This heuristic prioritizes the conditions based on the number of times they are used within the loop. The heuristic can be controlled with the option -unswitch-identical-branches-min-count=<n>.

Enables unswitching of a loop with respect to a branch conditional value (B), where B appears in at least <**n**> compares in the loop. This option is enabled with -aggressive-loop-unswitch. The default value is 3.

 ${f Usage:}$ -mllvm -aggressive-loop-unswitch -mllvm -unswitch-identical-branches-min-count=<n>

Where, **n** is a positive integer and lower value of <**n**> facilitates more unswitching.

-lv-function-specialization

Generates specialized function versions when the loops inside function are vectorizable and the arguments are not aliased with each other.

Usage: -mllvm -lv-function-specialization

• -loop-splitting

Enables splitting of loops into multiple loops to eliminate the branches, which compare the loop induction with an invariant or constant expression. By default, this option is enabled under -O3. To disable this optimization, use -loop-splitting=false.

Usage: -mllvm -loop-splitting

-enable-ipo-loop-split

Enables splitting of loops into multiple loops to eliminate the branches, which compares the loop induction with a constant expression. This constant expression can be derived through inter-

procedural analysis. By default, this option is enabled under -O3. To disable this optimization, use -enable-ipo-loop-split=false.

Usage: -mllvm -enable-ipo-loop-split

• -enable-loop-fusion

This option enables the classical loop fusion transformation where the bodies of multiple loop nests are fused into one loop nest. The transformation checks various legality criteria involving the bounds of the loop nests involved, the control flow nesting of the loop nests and so on. The transformation is off by default and may be enabled by the user by using this option. Loop fusion enables reuse of memory access operations across the loop nests and is also beneficial for cache performance. As part of the profitability check for this transformation it uses code size thresholds which control the size of the fused loop body created.

Usage: -mllvm -enable-loop-fusion

-enable-loopinterchange

This option enables the classical loop interchange or loop permutation transformation on a loop nest. It reorders the loops in a multi-dimensional loop nest, checking for various legality criteria in the process. The transformation is off by default and may be enabled by using this option. Loop interchange tries to find a reordering of the loops in a multi-dimensional loop nest such that the number of loop invariant expressions that may be hoisted out from an inner loop to a loop at a higher level may be maximized.

Usage: -mllvm -enable-loopinterchange

-fuse-tile-inner-loop

Enables fusion of adjacent tiled loops as a part of loop tiling transformation. This option is set to **false** by default.

Usage: -mllvm -fuse-tile-inner-loop

-enable-loop-distribute-adv

Enables advanced loop distribution which improves loop vectorization by separating out the portions of the loop affecting vectorization. This flag is disabled by default.

Usage: -mllvm -enable-loop-distribute-adv

-floop-transform

Enables the subset of loop transformations including improved variants of loop-fusion, loop-interchange, loop blocking, and distribution. It is enabled under -zopt.

Usage: -floop-transform

-faggressive-loop-transform

Enables the subset of loop transformations including improved variants of loop unswitching, loop-tiling, and loop versioning licm. It is enabled under -zopt.

Usage: -faggressive-loop-transform

-fstrip-mining

Enables loop strip mining optimization. This optimization breaks a large loop into smaller segments or strips to improve temporal and spatial locality. This option should be invoked with - flto.

```
Usage: -flto -fstrip-mining
```

-ldist-scalar-expand

Enables loop distribution with scalar expansion for better vectorization and it can be invoked as wl,-mllvm -Wl,-ldist-scalar-expand.

```
{f Usage:} -Wl,-mllvm -Wl,-ldist-scalar-expand
```

Math Options

```
-convert-pow-exp-to-int={true,false}
```

Converts the call to floating point exponent version of pow to its integer exponent version if the floating-point exponent can be converted to integer. This option is set to **true** by default.

```
Usage: -mllvm -convert-pow-exp-to-int={true, false}
```

Inline Optimization Options

-finline-aggressive

Enables improved inlining capability through better heuristics. This optimization is more effective when using with -flto as the whole program analysis is required to perform this optimization, which can be invoked as -flto -finline-aggressive.

```
Usage: -finline-aggressive
```

-inline-recursion=[1,2,3,4]

Enables inlining for recursive functions based on heuristics with level 4 being most aggressive. The default level will be 2. Higher levels may lead to code-bloat due to expansion of recursive functions at call sites.

- For level 1-2: Enables inlining for recursive functions using heuristics with inline depth 1. Level
 2 uses more aggressive heuristics.
- For level 3: Enables inlining for all recursive functions with inline depth 1.
- For level 4: Enables inlining for all recursive function with inline depth 10.

This is more effective with -flto as the whole program analysis is required to perform this optimization.

```
Usage: -flto -finline-recursion=[1,2,3,4]
```

Memory Layout Optimization Options

• -fstruct-layout={1,2,3,4,5,6,7,8,9}

Analyzes the whole program to determine if the structures in the code can be peeled, if dead or redundant fields can be deleted, and if the pointer or integer fields in the structure can be

compressed. If feasible, this optimization transforms the code to enable these improvements. This transformation is likely to improve cache utilization and memory bandwidth. It is expected to improve the scalability of programs executed on multiple cores.

This is effective only under -flto as the whole program analysis is required to perform this optimization. You can choose different levels of aggressiveness with which this optimization can be applied to your application; with 1 being the least aggressive and 7 being the most aggressive level.

- -fstruct-layout=0 disables structure peeling (default).
- -fstruct-layout=1 enables structure peeling.
- -fstruct-layout=2 enables structure peeling and selectively compresses self-referential pointers in these structures to 32-bit pointers, wherever safe.
- -fstruct-layout=3 enables structure peeling and selectively compresses self-referential pointers in these structures to 16-bit pointers, wherever safe.
- -fstruct-layout=4 enables structure peeling, pointer compression as in level 2 and further enables compression of structure fields that are of 64-bit to 32-bit integer type. This is performed under a strict safety check.
- -fstruct-layout=5 enables structure peeling, pointer compression as in level 3 and further enables compression of structure fields that are of 64-bit to 32-bit integer type. This is performed under a strict safety check.
- -fstruct-layout=6 enables structure peeling, pointer compression as in level 2 and further enables compression of structure fields that are of 64-bit to 32-bit integer type. This is performed under a strict safety check.
- -fstruct-layout=7 enables structure peeling, pointer compression as in level 3 and further enables compression of structure fields that are of 64-bit to 16-bit integer type.
- -fstruct-layout=8 enables structure peeling, pointer compression, 64-bit integer type compression as in level 6 and creates optimal ordering of peeled structure fields that could improve runtime performance.
- -fstruct-layout=9 enables structure peeling, pointer compression, 64-bit integer type compression as in level 7 and creates optimal ordering of peeled structure fields that could improve runtime performance.

Usage: -flto -fstruct-layout={1,2,3,4,5,6,7,8,9}

Notes:

- 1. -fstruct-layout=4 and -fstruct-layout=5 are derived from -fstruct-layout=2 and -fstruct-layout=3 respectively, with the added feature of safe compression of 64-bit to 32-bit integer fields in structures. Going from -fstruct-layout=4 to -fstruct-layout=5 may result in higher performance if the pointer values are such that the pointers can be compressed to 16-bits.
- 2. -fstruct-layout=6 and -fstruct-layout=7 are derived from -fstruct-layout=2 and -fstruct-layout=3 respectively, with the added feature of compression of 64-bit to 32-bit integer fields in structures. These are similar to -fstruct-layout=4 and -fstruct-layout=5, but here, the

integer fields of the structures are always compressed from 64-bits to 32-bits, without any safety guarantee.

• -fremap-arrays

Transforms the data layout of a single dimensional array to provide better cache locality. This optimization is effective only under -flto as the whole program analysis is required to perform this optimization, which can be invoked as -flto -fremap-arrays.

Usage: -flto -fremap-arrays

-reduce-array-computations={1,2,3}

Performs array dataflow analysis and optimizes the unused array computations.

- -reduce-array-computations=1: Eliminates the computations on unused array elements.
- -reduce-array-computations=2: Eliminates the computations on zero valued array elements.
- -reduce-array-computations=3: Eliminates the computations on unused and zero valued array elements (combination of 1 and 2).

This optimization is effective with -flto as the whole program analysis is required to perform this optimization, which can be invoked as -flto -reduce-array-computations={1,2,3}.

Usage: -flto -mllvm -reduce-array-computations={1,2,3}

Instruction Level Optimization Options

-enable-X86-prefetching

Enables the generation of x86 prefetch instruction for the memory references inside a loop/ inside an inner most loop of a loop nest to prefetch the second dimension of multidimensional array/ memory references in the inner most of a loop nest.

Usage: -mllvm -enable-X86-prefetching

-suppress-fmas

Identifies the reduction patterns on FMA and suppresses the FMA generation as it is not profitable on the reduction patterns.

Usage: -mllvm -suppress-fmas

-fnt-store

Generates a non-temporal store instruction for array accesses in a loop with a large trip count.

Usage: -fnt-store

-fnt-store=aggressive

Generates non-temporal store instruction for array accesses in a loop, whose iteration count cannot be determined at compile time. In this case, compiler assumes the iteration count is huge.

Usage: -fnt-store=aggressive

-enable-redundant-movs

Removes any redundant mov operations including redundant loads from memory and stores to memory. This can be invoked using -Wl,-plugin-opt=-enable-redundant-movs.

Usage: -Wl,-plugin-opt=-enable-redundant-movs

• -merge-constant

Attempts to promote frequently occurring constants to registers. The aim is to reduce the size of the instruction encoding for instructions using constants and obtain a performance improvement.

Usage: -mllvm -merge-constant

-optimize-strided-mem-cost

Optimizes the cost model for strided access to memory.

Usage: -mllvm -optimize-strided-mem-cost

-fenable-aggressive-gather

This option enables generation of gather instructions for the cases where it is profitable.

Usage: -fenable-aggressive-gather

Scalar Optimization Options

-enable-licm-vrp

Enables estimation of the virtual register pressure before performing loop invariant code motion. This estimation is used to control the number of loop invariants that will be hoisted during the loop invariant code motion.

Usage: -mllvm -enable-licm-vrp

-do-block-reorder={none, simple, aggressive}

Reorders the control predicates in increasing order of complexity from outer predicate to inner. This option is set to simple by default. The simple mode reorders simple expressions while the **aggressive** mode reorders the predicates involving function calls despite the presence of code dealing with exceptions.

This optimization also includes safety analysis which checks if it is safe to reorder the basic blocks. However, when this optimization takes effect, the safety analysis ignores exceptions specifically within the call chain involved in the context of the blocks being reordered. This would be acceptable for most input programs, but if your program strictly needs to support throwing of exceptions at all points in the program, it is advisable to avoid using this option.

Usage: -mllvm -do-block-reorder={none, simple, aggressive}

-fscalar-transform

Enables a subset of scalar transformations including improved variants of various code movement optimizations such as hoisting and invariant code motion. It is enabled under -zopt.

Usage: -fscalar-transform

Miscellaneous Options

-fitodcalls

Promotes indirect to direct calls by placing conditional calls. Application or benchmarks that have small and deterministic set of target functions for function pointers that are passed as call parameters benefit from this optimization. Indirect-to-direct call promotion transforms the code to use all possible determined targets under runtime checks and falls back to the original code for all the other cases. Runtime checks are introduced by the compiler for each of these possible function pointer targets followed by direct calls to the targets.

This is a link time optimization, which is invoked as -flto -fitodcalls.

Usage: -flto -fitodcalls

-fitodcallsbyclone

Performs value specialization for functions with function pointers passed as an argument. It does this specialization by generating a clone of the function. The cloning of the function happens in the call chain as needed to allow conversion of indirect function call to direct call. This complements -fitodcalls optimization and is also a link time optimization, which is invoked as -flto -fitodcallsbyclone.

Usage: -flto -fitodcallsbyclone

-function-specialize

Optimizes the functions with compile time constant formal arguments. This option is enabled by default under -O3.

Usage: -mllvm -function-specialize

• -favoid-fpe-causing-opt

Restricts a few optimizations that leads to floating point exceptions.

Usage: -favoid-fpe-causing-opt

Diagnostics Option(s)

-fsanitize

Runs the sanitizers for diagnostics.

Usage: -fsanitize={address,thread,memory,safe-stack}

5.4 Deprecated Options

The following options have been deprecated:

- -vectorize-memory-aggressively (from AOCC 2.2.0)
- -Menable-vectorize-pragmas=<value> (from AOCC 3.2.0)

Chapter 6 Debuggability

6.1 OpenMP Debugging Support (OMPD)

Note: This is available in AOCC 2.3 or later.

The AOCC installation includes OMPD for debugging C/C++ OpenMP programs through a gdb plugin with limited functionality.

Note: Debugging the code that runs on an offloading device is not supported.

Complete the following steps to use OMPD for debugging C/C++ OpenMP programs through a gdb plugin:

Note: For using the OMPD plugin, Python 3.5 or later is required.

1. Add folders **ompd** and **lib** to your **LD LIBRARY PATH** using this command:

\$ export LD_LIBRARY_PATH=<compdir>/aocc-compiler-<ver>/ompd:<compdir>/aocc-compiler-<ver>/
lib:\$LD_LIBRARY_PATH

2. Set OMP DEBUG to enabled:

\$ export OMP_DEBUG=enabled

3. Compile the program to be debugged with **-g** and **-fopenmp** options as follows for a sample C source file *xyz.c*:

```
$ <compdir>/aocc-compiler-<ver>/bin/clang -g -fopenmp xyz.c -o xyz.out
```

Note: The program to be debugged needs to have a dynamic link dependency on 'libomp.so' under <compdir>/aocc-compiler-<ver>/lib for OpenMP-specific debugging to work correctly. The user can check this using ldd on the generated binary, that is xyz.out.

4. Debug the binary xyz.out by invoking gdb with the plugin as follows:

```
$ gdb -x <compdir>/aocc-compiler-<ver>/ompd/__init__.py ./xyz.out
```

Note: The plugin <compdir>/aocc-compiler-<ver>/ompd/ init .py must be used.

6.2 OMPD Commands

The following table describes the OMPD commands:

Table 2. OMPD Commands

| Command | Description |
|-----------|---|
| help ompd | It lists the subcommands available for OpenMP specific debugging. |

Table 2. OMPD Commands

| Command | Description |
|-----------|---|
| ompd init | It must be run first to load the libompd.so available in the \$LD_LIBRARY_PATH environment variable and to initialize the OMPD library. It starts the program run and the program stops at a temporary breakpoint at the OpenMP internal location ompd_dll_locations_valid(). You can continue from the temporary breakpoint for debugging. You can place breakpoints at the OpenMP internal locations ompd_bp_thread_begin and ompd_bp_thread_end to catch the begin and end events ompd_bp_task_begin and ompd_bp_task_end breakpoints can be used to catch the beginning and ending of the events ompd_bp_parallel_begin and ompd_bp_parallel_end can be used to catch the beginning and ending of the parallel events. |

6.3 OMPD Subcommands

The following table lists the OMPD subcommands that can used inside gdb:

Table 3. OMPD Subcommands

| Subcommand | Description |
|---------------|--|
| ompd init | Finds and initializes the OMPD library. |
| ompd bt | Used to turn the filter <i>on</i> or <i>off</i> for the bt output on or off. You must specify the <i>on continued</i> option to trace the worker threads back to the master threads. |
| ompd icvs | Displays the values of the Internal Control Variables. |
| ompd parallel | Displays the details of the current and enclosing parallel regions. |
| ompd step | Executes step and skip runtime frames as much as possible. |
| ompd threads | Provides the details of the current threads. |

Chapter 7 Diagnostics

7.1 **AOCC Optimization Report (AOR)**

The AOR tool generates an optimization report for any C, C++, and Fortran application project compiled with AOCC. The optimization report can be investigated by the application developer for a list of optimizations performed/not performed on the compiled application with a hint for the course of action. Currently, AOR contains the report of messages related to loop-vectorize optimization pass information and can be generated with the following commands:

• -fgen-aor

For the compile-time AOR generation (CFLAGS, CXXFLAGS, and FCFLAG).

-fgen-aor -flto -fuse-ld=lld

For the link-time AOR generation (CFLAGS, CXXFLAGS, and FCFLAG).

Example:

```
git clone https://github.com/UoB-HPC/TSVC_2
cd TSVC_2/src
clang -03 -march=znver2 -fgen-aor -flto -fuse-ld=lld *.c -I . -o tsvc -lm
```

The above commands will generate a compile-time report *optimization_report.aor* and a link-time report *optimization_ld report.aor* respectively in the root folder of the application project.

Prerequisites

AOR requires the following packages to be installed:

- Python 3.8
- Python-YAML
- Python3-pip
- pip3 packages including Pygments and PyYAML

Chapter 8 Support

For support options, the latest documentation, and downloads refer AMD Developer Central (https://developer.amd.com/amd-aocc/).

Chapter 9 References

The following document has been used as a reference for this document:

LLVM Documentation (https://releases.llvm.org/14.0.0/tools/clang/docs/ClangCommandLineReference.html)